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10/733,962	12/11/2003	Terrenc J. Tanis	PB 06 0006 (SPLG 2)	8771
38790 7590 07/05/2007 THE SMALL PATENT LAW GROUP LLP 611 OLIVE STREET, SUITE1611 ST. LOUIS, MO 63101			EXAMINER KAO, JUTAI	
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Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

<b>Office Action Summary</b>	<b>Application No.</b> 10/733,962	<b>Applicant(s)</b> TANIS ET AL.	
	<b>Examiner</b> Ju-Tai Kao	<b>Art Unit</b> 2616	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

#### Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

#### Status

- 1) ☐ Responsive to communication(s) filed on \_\_\_\_.
- 2a) ☐ This action is **FINAL**.                      2b) ☒ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

#### Disposition of Claims

- 4) ☒ Claim(s) 1-25 is/are pending in the application.
- 4a) Of the above claim(s) \_\_\_\_ is/are withdrawn from consideration.
- 5) ☐ Claim(s) \_\_\_\_ is/are allowed.
- 6) ☒ Claim(s) 1-25 is/are rejected.
- 7) ☐ Claim(s) \_\_\_\_ is/are objected to.
- 8) ☐ Claim(s) \_\_\_\_ are subject to restriction and/or election requirement.

#### Application Papers

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☒ The drawing(s) filed on 11 December 2003 is/are: a) ☒ accepted or b) ☐ objected to by the Examiner.  
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).  
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

#### Priority under 35 U.S.C. § 119

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All    b) ☐ Some \* c) ☐ None of:
1. ☐ Certified copies of the priority documents have been received.
  2. ☐ Certified copies of the priority documents have been received in Application No. \_\_\_\_.
  3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

\* See the attached detailed Office action for a list of the certified copies not received.

#### Attachment(s)

- |   |  |
|---|--|
| 1) <input checked="" type="checkbox"/> Notice of References Cited (PTO-892)   | 4) <input type="checkbox"/> Interview Summary (PTO-413)<br>Paper No(s)/Mail Date. ____ |
| 2) <input type="checkbox"/> Notice of Draftsperson's Patent Drawing Review (PTO-948)  | 5) <input type="checkbox"/> Notice of Informal Patent Application                      |
| 3) <input checked="" type="checkbox"/> Information Disclosure Statement(s) (PTO/SB/08)<br>Paper No(s)/Mail Date <u>12/11/2003</u> . | 6) <input type="checkbox"/> Other: ____  |

## **DETAILED ACTION**

### ***Claim Rejections - 35 USC § 112***

1. The following is a quotation of the second paragraph of 35 U.S.C. 112:

The specification shall conclude with one or more claims particularly pointing out and distinctly claiming the subject matter which the applicant regards as his invention.

2. Claim 1-25 are rejected under 35 U.S.C. 112, second paragraph, as being indefinite for failing to particularly point out and distinctly claim the subject matter which applicant regards as the invention.

Regarding claim 1-25, random variables "p", "nth", "m", and "mth" are recited without being defined.

Regarding claim 9, line 11 of the claim recites "the p timeslots". However, it is unclear which set of p timeslots is being referred to as the parent claim recites "a first set of p timeslots" and claim 9 also recites a separate "p timeslots of a second route". In addition, claim 9, recites "a uniform search". However, it is unclear what a "uniform search" is. For the purpose of examining this claim, a "uniform search" is considered as a method of searching.

Regarding claim 12, "a uniform search" is recited. However, it is unclear what a "uniform search" is. For the purpose of examining this claim, a "uniform search" is considered as a method of searching.

Regarding claim 24, "a first, a second, and a third data collection" is recited in line 3. However, the parent claim already recites "a first data collection". Therefore, it is unclear whether "a first...data collection" of claim 24 refers to the same data collection

recited in its parent claim or a separate data collection. For the purpose of examination, the two "a first data collection" are considered as the same data collection.

***Claim Rejections - 35 USC § 102***

3. The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless –

(e) the invention was described in a patent granted on an application for patent by another filed in the United States before the invention thereof by the applicant for patent, or on an international application by another who has fulfilled the requirements of paragraphs (1), (2), and (4) of section 371(c) of this title before the invention thereof by the applicant for patent.

The changes made to 35 U.S.C. 102(e) by the American Inventors Protection Act of 1999 (AIPA) and the Intellectual Property and High Technology Technical Amendments Act of 2002 do not apply when the reference is a U.S. patent resulting directly or indirectly from an international application filed before November 29, 2000. Therefore, the prior art date of the reference is determined under 35 U.S.C. 102(e) prior to the amendment by the AIPA (pre-AIPA 35 U.S.C. 102(e)).

4. Claim 1-5, and 17 are rejected under 35 U.S.C. 102(e) as being anticipated by Williams (US 6,356,550).

Williams discloses a flexible time division multiplexed bus using SONET formatting including the following features.

Regarding claim 1, a method for switching multi-rate communications (see "communications system transports facility signals including: DS1, DS2, and DS3

signals, a 2.048 Mbps E1 signal...in addition to the 51.84 Mbps SONET STS-1 level signal" recited in column 6, line 63-67) comprising: filling a first set of  $p$  timeslots from a first data collection of a first type (see Fig. 2J, timeslot 3 and 18, the two timeslots filled with 2 sets ( $p = 2$ ) of data from data collection M of the Virtual Tributary 3 type (VT3); Fig. 2H shows that collection M is a VT3), wherein the first data collection includes  $m$  sets of data and  $p$  is less than  $m$  (see Fig. 2I, where 6 sets of M slots are shown, and the number of M slots in the VT3 is greater than 2, that is,  $m = 6$ ); and loading at least one overflow timeslot (see Fig. 2J, slot 33) with at least one overflow set of data from the first data collection (the slot being filled with overflowing M slots from the 6 VT3 slots shown in Fig. 2H), wherein the overflow set of data exceeds a size of the first set of  $p$  timeslots (Fig. 2H, besides the first two M slots filled into Fig. 2J described above, there are 4 more overflowing M slots in Fig. 2H).

Regarding claim 2, filling a second set of  $p$  timeslots with  $p$  sets of data from a second data collection of the first type (see Fig. 2J, slot 11 and 25, the two timeslots filled with two sets of N data from the data collection N of the VT3 type as shown in Fig. 2H), wherein the second data collection includes  $m$  sets of data (see Fig. 2H, data collection N also includes 6 slots of data).

Regarding claim 3, loading the at least one overflow timeslot (see Fig. 2J, slot 11, overflowing since it's not one of the first set of  $p$  timeslots) with at least one overflow set of data (see Fig. 2J, slot 11, which is filled with data N) from a second data collection (Fig. 2J, slot 11 is filled with data from the N data collection) of the first type (see Fig.

2H, N is also of the VT3 type), wherein the second data collection includes m sets of data (see Fig. 2H, data collection N also includes 6 sets of data).

Regarding claim 4, filling a second set of p timeslots with p sets of data from a second data collection (see Fig. 2J, slots 4 and 12, which are 2 sets of timeslots filled with 2 sets of data from the data collection O), wherein the second data collection is of a second type that is different than the first type (see Fig. 2I, where data collection O is of the VT6 type) and the second data collection includes p sets of data (see Fig. 2I, data collection includes 2 sets of data, and a total of 12 sets of data).

Regarding claim 5, filling a second set of p timeslots with p sets of data from a second data collection (see Fig. 2J, two sets of timeslots 1 and 31, each filled with two sets of data A from the data collection A), wherein the second data collection is of a second type that is different than the first type (see Fig. 2F, where data collection A is of the type VT1), wherein the second type is a DS1 type (see "to transmit DS1 signal, VT1 uses..." recited in column 2, line 36); and filling a second set of p timeslots with p sets of data from the second data collection (see Fig. 2J, two sets of timeslots 1 and 31, each filled with two sets of data A from the data collection A).

Regarding claim 17, a method for switching multi-rate communications (see "communications system transports facility signals including: DS1, DS2, and DS3 signals, a 2.048 Mbps E1 signal...in addition to the 51.84 Mbps SONET STS-1 level signal" recited in column 6, line 63-67) comprising: filling a first set of p timeslots (see Fig. 2J-2K, timeslot 2, 24 and 46, that is,  $p = 3$ ) with p sets of data of a first data collection (see Fig. 2J, timeslots 2, 24, and 46 are filled with three sets of data from the

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data collection X, which is shown in Fig. 2G), a second set of p timeslots (see Fig. 2J-K, timeslots 10, 32, and 53) with p sets of data of a second data collection (see Fig. 2J, timeslot 10, 32, and 53 are filled with three sets of data from the data collection Y, which is shown in Fig. 2G), a third set of p timeslots (see Fig. 2J, timeslot 17, 37, and 61) with p sets of data of a second data collection (see Fig. 2J, timeslot 17, 37, and 61 are filled with three sets of data from the data collection Z, which is shown in Fig. 2G), wherein each of the first, the second, and the third data collection are of a common first type (see Fig. 2G, showing that each of data collection X, Y and Z are of the type VT2), each of the first, the second, and the third data collection include m sets of data, and p is less than m (see Fig. 2G, where each of the VT2 include 4 sets of data, that is  $m = 4$ ); and loading a fourth set of p timeslots (see Fig. 2J, timeslots 68, 77, and 82) with an mth set of data of the first data collection, an mth set of data of the second data collection, and an mth set of data of the third data collection (see Fig. 2J, where timeslots 68, 77, and 82 are filled with the fourth set of data X, Y and Z respectively).

### ***Claim Rejections - 35 USC § 103***

5. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

6. This application currently names joint inventors. In considering patentability of the claims under 35 U.S.C. 103(a), the examiner presumes that the subject matter of

the various claims was commonly owned at the time any inventions covered therein were made absent any evidence to the contrary. Applicant is advised of the obligation under 37 CFR 1.56 to point out the inventor and invention dates of each claim that was not commonly owned at the time a later invention was made in order for the examiner to consider the applicability of 35 U.S.C. 103(c) and potential 35 U.S.C. 102(e), (f) or (g) prior art under 35 U.S.C. 103(a).

7. Claim 6 is rejected under 35 U.S.C. 103(a) as being unpatentable over Williams (US 6,356,550) in view of Eyeson (US 2003/0223568) and Boily (US 2004/0001454).

Williams discloses the claimed limitations above.

Williams does not disclose the following features: regarding claim 6, wherein a network system includes a number of timeslot interchanges, said method further comprising determining whether a capacity of the timeslot interchanges to support a number of connections of the first type and a second type is exceeded if a connection is routed via the timeslot interchanges.

Eyeson discloses a method and apparatus for automatically directing calls by an invisible agent in a switch including the following features.

Regarding claim 6, wherein a network system (see Fig. 1, which shows a WAN network) includes a number of timeslot interchanges (see Fig. 1, which includes remote switch 102 and 108, and Fig. 2, which shows the detail of each remote switch and includes time slot interchange TSI 206).

Boily discloses a timeslot interchange switch including the following features.



Regarding claim 6, said method further comprising determining whether a capacity of the timeslot interchanges to support a number of connections (see “there is an upper limit to the number of connections that system 10 is capable of processing...system 10 may maintain a count of active locations 48 and signal an alarm if a number of active locations 48 exceeds a threshold” recited in paragraph 43, wherein the system 10 is part of a time slot interchange switch shown in Fig. 1, and active location represents connections) of the first type (described in the rejection to claim 1 above) and a second type (described in Williams, e.g. data type N in Fig. 2J of Williams) is exceeded if a connection is routed via the timeslot interchanges (again, system 10 being the timeslot interchange switch).

It would have been obvious for one of the ordinary skill in the art at the time of the invention to modify the system of Williams using features, as taught by Eyeson and Boily, in order to provide a non-blocking multi-rate transmission in a communication network involving different type of signals.

8. Claim 7, 10, and 18 are rejected under 35 U.S.C. 103(a) as being unpatentable over Williams (US 6,356,550) in view of Aicklen (US 7,145,867) and Eyeson (US 2003/0223568).

Williams discloses the claimed limitations above. Williams does not disclose the following features: regarding claim 7, wherein a network system includes at least two timeslot interchanges coupled to at least one space switch, different first and second routes exist via the at least two timeslot interchanges and the at least one space switch,

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said method further comprising: routing the p sets of data of the first data collection via the first route if the first route is unblocked; and selecting p timeslots of the second route if the first route is blocked; regarding claim 10, wherein a network system includes at least two timeslot interchanges coupled to at least one space switch, different first and second routes exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising selecting an nth timeslot of the second route if the first route is blocked; regarding claim 18, wherein a network system includes at least two timeslot interchanges coupled to at least one space switch, a first route and a second route exists via the at least two timeslot interchanges and the at least one space switch, and the first route is different than the second route, said method further comprising: routing the p sets of data of the first data collection via the first route if the first route is unblocked; and selecting p timeslots of the second route different than the first route if the first route is blocked

Aicklen discloses the system and method for slot deflection routing including the following features.

Regarding claim 7, wherein a network system (see network system including five edge units in Fig. 7) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 7, edge units 0-4; each edge unit could can be a router or switch, see "router...edge unit" recited in column 1, line 32-35; and each router could contain a timeslot interchange and a space switch as will be explained below using the Eyeson reference), different first and second routes (see Fig. 7, route 335 and route 333-2, 333-5) exist via the at least two timeslot interchanges and the at least one space switch, said

method further comprising: routing the p sets of data of the first data collection (explained in the rejection made to claim 1) via the first route if the first route is unblocked (see Fig. 6, where the route between edge unit 2 and 3 are unblocked, and data in timeslot T0 is transferred in route 333-0); and selecting p timeslots of the second route if the first route is blocked (see Fig. 7, where the route between edge unit 2 and 3 are blocked, and data are transferred through 333-2, 333-5 in timeslot TS2 and TS5; also see "a failure in the link..." recited in column 11, line 37-55).

Regarding claim 10, wherein a network system (see network system including five edge units in Fig. 7) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 7, edge units 0-4; each edge unit could can be a router or switch, see "router...edge unit" recited in column 1, line 32-35; and each router could contain a timeslot interchange and a space switch as will be explained below using the Eyeson reference), different first and second routes (see Fig. 7, route 335 and route 333-2, 333-5) exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising selecting an nth timeslot (see timeslot T2, in Fig. 7) of the second route (see route 333-2 in Fig. 7) if the first route is blocked (see "a failure in the link from EU2 to EU3..." recited in column 11, line 37-55).

Regarding claim 18, wherein a network system (see network system including five edge units in Fig. 7) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 7, edge units 0-4; each edge unit could can be a router or switch, see "router...edge unit" recited in column 1, line 32-35; and each router could contain a timeslot interchange and a space switch as will be explained below using the

Eyeson reference), different first and second routes (see Fig. 7, route 335 and route 333-2, 333-5) exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising: routing the p sets of data of the first data collection (explained above in the rejection made in claim 17) via the first route if the first route is unblocked (see Fig. 6, where the route between edge unit 2 and 3 are unblocked, and data in timeslot T0 is transferred in route 333-0); and selecting p timeslots of the second route different than the first route if the first route is blocked (see Fig. 7, where the route between edge unit 2 and 3 are blocked, and data are transferred through 333-2, 333-5 in timeslot TS2 and TS5; also see "a failure in the link..." recited in column 11, line 37-55).

Eyeson discloses the method and apparatus for automatically directing calls by an invisible agent in a switch.

Regarding claim 7, wherein a network system (explained in Aicklen above) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 1 showing remote switch 102, which is again shown in Fig. 2 showing that the remote switch includes a router portion 221 and a circuit switch portion 223; the router portion includes a time slot interchange portion TSI 206, and the circuit switch being a space switch; thus, each of the edge unit in Aicklen can be modified to include a TSI coupled to a router, thus each route in Aicklen includes two TSI coupled to two space switch from the two edge units connected to form the route), different first and second routes exist via the at least two timeslot interchanges and the at least one space

switches (explained above, each routes in Aicklen includes two TSI coupled to two space switches).

Regarding claim 10, wherein a network system (explained in Aicklen above) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 1 showing remote switch 102, which is again shown in Fig. 2 showing that the remote switch includes a router portion 221 and a circuit switch portion 223; the router portion includes a time slot interchange portion TSI 206, and the circuit switch being a space switch; thus, each of the edge unit in Aicklen can be modified to include a TSI coupled to a router, thus each route in Aicklen includes two TSI coupled to two space switch from the two edge units connected to form the route), different first and second routes exist via the at least two timeslot interchanges and the at least one space switches (explained above, each routes in Aicklen includes two TSI coupled to two space switches).

Regarding claim 18, wherein a network system (explained in Aicklen above) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 1 showing remote switch 102, which is again shown in Fig. 2 showing that the remote switch includes a router portion 221 and a circuit switch portion 223; the router portion includes a time slot interchange portion TSI 206, and the circuit switch being a space switch; thus, each of the edge unit in Aicklen can be modified to include a TSI coupled to a router, thus each route in Aicklen includes two TSI coupled to two space switch from the two edge units connected to form the route), different first and second routes exist via the at least two timeslot interchanges and the at least one space

switches (explained above, each routes in Aicklen includes two TSI coupled to two space switches).

It would have been obvious for one of the ordinary skill in the art at the time of the invention to modify the system of Williams using features, as taught by Aicklen and Eyeson, in order to provide a non-blocking multi-rate transmission in a communication network involving different type of signals and prevent packets being dropped from path failure.

9. Claim 8, 9, 11, 12, and 19-20 are rejected under 35 U.S.C. 103(a) as being unpatentable over Williams (US 6,356,550) in view of Aicklen (US 7,145,867) Eyeson (US 2003/0223568) and Akahane (US 2001/0050914).

Williams discloses the claimed limitations above. Williams does not disclose the following features: regarding claim 8, wherein a network system includes at least two timeslot interchanges coupled to at least one space switch, different first and second routes exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising: selecting p timeslots of the second route if the first route is blocked; and performing a sequential search for the second route by identifying, in a consecutive fashion within the at least two timeslot interchanges, the p timeslots; regarding claim 9, wherein a network system includes at least two timeslot interchanges coupled to at least one space switch, a first route and a second route exists via the at least two timeslot interchanges and the at least one space switch, and the first route is different than the second route, said method further comprising: selecting p timeslots of

a second route different than a first route via at least two timeslot interchanges and at least one space switch for transporting the  $p$  sets of data from the first data collection if the first route is blocked; and performing a uniform search for the second route by determining whether the  $p$  timeslots have less load than loads of remaining timeslots within the at least two timeslot interchanges; regarding claim 11, wherein a network system includes at least two timeslot interchanges coupled to at least one space switch, different first route and second routes exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising: selecting an  $n$ th timeslot of the second route if the first route is blocked; and performing a sequential search for the second route by identifying, in a consecutive fashion within the at least two timeslot interchanges, the  $n$ th timeslot; regarding claim 12, wherein a network system includes at least two timeslot interchanges coupled to at least one space switch, different first route and second routes exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising: selecting an  $n$ th timeslot of the second route if the first route is blocked; and performing a uniform search for the second route by determining whether the  $n$ th timeslot has less load than loads of remaining timeslots within the at least two timeslot interchanges; regarding claim 19, wherein a network system includes at least two timeslot interchanges coupled to at least one space switch, a first route and a second route exists via the at least two timeslot interchanges and the at least one space switch, and the first route is different than the second route, said method further comprising: selecting  $p$  timeslots within the second route for transporting the  $p$  sets of data from the first data collection if the first route is

blocked; and performing a sequential search for the second route by identifying, in a consecutive fashion within the at least two timeslot interchanges, the p timeslots within the second route; regarding claim 20, wherein a network system includes at least two timeslot interchanges coupled to at least one space switch, a first route and a second route exists via the at least two timeslot interchanges and the at least one space switch, and the first route is different than the second route, said method further comprising: selecting an nth timeslot of the second route if the first route is blocked; and performing a uniform search for the second route by determining whether the nth timeslot has less load than loads of remaining timeslots within the at least two timeslot interchanges.

Aicklen discloses the system and method for slot deflection routing including the following features.

Regarding claim 8, wherein a network system (see network system including five edge units in Fig. 7) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 7, edge units 0-4; each edge unit could can be a router or switch, see "router...edge unit" recited in column 1, line 32-35; and each router could contain a timeslot interchange and a space switch as will be explained below using the Eyeson reference), different first and second routes (see Fig. 7, route 335 and route 333-2, 333-5) exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising: selecting p timeslots of the second route (see timeslot T2 and T5 in Fig. 7) if the first route is blocked (see "a failure in the link from EU2 to EU3..." recited in column 11, line 37-55).



Regarding claim 9, wherein a network system (see network system including five edge units in Fig. 7) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 7, edge units 0-4; each edge unit could can be a router or switch, see "router...edge unit" recited in column 1, line 32-35; and each router could contain a timeslot interchange and a space switch as will be explained below using the Eyeson reference), different first and second routes (see Fig. 7, route 335 and route 333-2, 333-5) exist via the at least two timeslot interchanges and the at least one space switch and the first route is different than the second route (see Fig. 7, first route being route 335 and second route being route 333-2, 333-5), said method further comprising: selecting p timeslots of the second route (see route 333-2 and 333-5 carrying timeslot T2 and T5 in Fig. 7) different than a first route (see Fig. 7, route 335) via at least two timeslot interchanges and at least one space switch (see Fig. 7, second route includes edge units 0, 2 and 3, and each edge unit includes a TSI and a space switch as explained below in Eyeson) for transporting the p sets of data from the first data collection (explained in Williams above) if the first route is blocked (see "a failure in the link from EU2 to EU3..." recited in column 11, line 37-55); determining whether the p timeslots have less load than loads of remaining timeslots within the at least two timeslot interchanges (see "Now, consider a doubling in traffic..." recited in column 10, line 22-42, wherein the reference describes that the original route between two edge units could not handle the extra traffic, and the system looks to route the extra traffic through a second route with a lower capacity fill ratio that could handle the extra traffic).

Regarding claim 11, wherein a network system (see network system including five edge units in Fig. 7) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 7, edge units 0-4; each edge unit could can be a router or switch, see "router...edge unit" recited in column 1, line 32-35; and each router could contain a timeslot interchange and a space switch as will be explained below using the Eyeson reference), different first and second routes (see Fig. 7, route 335 and route 333-2, 333-5) exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising: selecting an nth timeslot of the second route (see timeslot T2 and T5 in Fig. 7) if the first route is blocked (see "a failure in the link from EU2 to EU3..." recited in column 11, line 37-55).

Regarding claim 12, wherein a network system (see network system including five edge units in Fig. 7) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 7, edge units 0-4; each edge unit could can be a router or switch, see "router...edge unit" recited in column 1, line 32-35; and each router could contain a timeslot interchange and a space switch as will be explained below using the Eyeson reference), different first and second routes (see Fig. 7, route 335 and route 333-2, 333-5) exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising: selecting an nth timeslot of the second route (see timeslot T2 and T5 in Fig. 7) if the first route is blocked (see "a failure in the link from EU2 to EU3..." recited in column 11, line 37-55); determining whether the nth timeslot has less load than loads of remaining timeslots within the at least two timeslot interchanges (see "Now, consider a doubling in traffic..." recited in column 10, line 22-

42, wherein the reference describes that the original route between two edge units could not handle the extra traffic, and the system looks to route the extra traffic through a second route with a lower capacity fill ratio that could handle the extra traffic).

Regarding claim 19, wherein a network system (see network system including five edge units in Fig. 7) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 7, edge units 0-4; each edge unit could can be a router or switch, see "router...edge unit" recited in column 1, line 32-35; and each router could contain a timeslot interchange and a space switch as will be explained below using the Eyeson reference), different first and second routes (see Fig. 7, route 335 and route 333-2, 333-5) exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising: selecting p timeslots within the second route for transporting the p sets of data (p sets of data are explained in the rejection to claim 17 above) from the first data collection if the first route is blocked (see Fig. 7, where the route between edge unit 2 and 3 are blocked, and data are transferred through 333-2, 333-5 in timeslot TS2 and TS5; also see "a failure in the link..." recited in column 11, line 37-55).

Regarding claim 20, wherein a network system (see network system including five edge units in Fig. 7) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 7, edge units 0-4; each edge unit could can be a router or switch, see "router...edge unit" recited in column 1, line 32-35; and each router could contain a timeslot interchange and a space switch as will be explained below using the Eyeson reference), different first and second routes (see Fig. 7, route 335 and route

333-2, 333-5) exist via the at least two timeslot interchanges and the at least one space switch, said method further comprising: selecting an nth timeslot of the second route (see timeslot T2 and T5 in Fig. 7) if the first route is blocked (see "a failure in the link from EU2 to EU3..." recited in column 11, line 37-55); determining whether the nth timeslot has less load than loads of remaining timeslots within the at least two timeslot interchanges (see "Now, consider a doubling in traffic..." recited in column 10, line 22-42, wherein the reference describes that the original route between two edge units could not handle the extra traffic, and the system looks to route the extra traffic through a second route with a lower capacity fill ratio that could handle the extra traffic).

Eyeson discloses the method and apparatus for automatically directing calls by an invisible agent in a switch.

Regarding claim 8, wherein a network system (explained in Aicklen above) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 1 showing remote switch 102, which is again shown in Fig. 2 showing that the remote switch includes a router portion 221 and a circuit switch portion 223; the router portion includes a time slot interchange portion TSI 206, and the circuit switch being a space switch; thus, each of the edge unit in Aicklen can be modified to include a TSI coupled to a router, thus each route in Aicklen includes two TSI coupled to two space switch from the two edge units connected to form the route), different first and second routes exist via the at least two timeslot interchanges and the at least one space switches (explained above, each routes in Aicklen includes two TSI coupled to two space switches).

Regarding claim 9, wherein a network system (explained in Aicklen above) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 1 showing remote switch 102, which is again shown in Fig. 2 showing that the remote switch includes a router portion 221 and a circuit switch portion 223; the router portion includes a time slot interchange portion TSI 206, and the circuit switch being a space switch; thus, each of the edge unit in Aicklen can be modified to include a TSI coupled to a router, thus each route in Aicklen includes two TSI coupled to two space switch from the two edge units connected to form the route), different first and second routes exist via the at least two timeslot interchanges and the at least one space switches (explained above, each routes in Aicklen includes two TSI coupled to two space switches).

Regarding claim 11, wherein a network system (explained in Aicklen above) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 1 showing remote switch 102, which is again shown in Fig. 2 showing that the remote switch includes a router portion 221 and a circuit switch portion 223; the router portion includes a time slot interchange portion TSI 206, and the circuit switch being a space switch; thus, each of the edge unit in Aicklen can be modified to include a TSI coupled to a router, thus each route in Aicklen includes two TSI coupled to two space switch from the two edge units connected to form the route), different first and second routes exist via the at least two timeslot interchanges and the at least one space switches (explained above, each routes in Aicklen includes two TSI coupled to two space switches).

Regarding claim 12, wherein a network system (explained in Aicklen above) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 1 showing remote switch 102, which is again shown in Fig. 2 showing that the remote switch includes a router portion 221 and a circuit switch portion 223; the router portion includes a time slot interchange portion TSI 206, and the circuit switch being a space switch; thus, each of the edge unit in Aicklen can be modified to include a TSI coupled to a router, thus each route in Aicklen includes two TSI coupled to two space switch from the two edge units connected to form the route), different first and second routes exist via the at least two timeslot interchanges and the at least one space switches (explained above, each routes in Aicklen includes two TSI coupled to two space switches).

Regarding claim 19, wherein a network system (explained in Aicklen above) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 1 showing remote switch 102, which is again shown in Fig. 2 showing that the remote switch includes a router portion 221 and a circuit switch portion 223; the router portion includes a time slot interchange portion TSI 206, and the circuit switch being a space switch; thus, each of the edge unit in Aicklen can be modified to include a TSI coupled to a router, thus each route in Aicklen includes two TSI coupled to two space switch from the two edge units connected to form the route), different first and second routes exist via the at least two timeslot interchanges and the at least one space switches (explained above, each routes in Aicklen includes two TSI coupled to two space switches).

Regarding claim 20, wherein a network system (explained in Aicklen above) includes at least two timeslot interchanges coupled to at least one space switch (see Fig. 1 showing remote switch 102, which is again shown in Fig. 2 showing that the remote switch includes a router portion 221 and a circuit switch portion 223; the router portion includes a time slot interchange portion TSI 206, and the circuit switch being a space switch; thus, each of the edge unit in Aicklen can be modified to include a TSI coupled to a router, thus each route in Aicklen includes two TSI coupled to two space switch from the two edge units connected to form the route), different first and second routes exist via the at least two timeslot interchanges and the at least one space switches (explained above, each routes in Aicklen includes two TSI coupled to two space switches).

Akahane discloses the VPN router and VPN identification method by using logical channel identifiers including the following features.

Regarding claim 8, performing a sequential search (see "a sequential table search method" recited in paragraph 72, where the tables are VPN identification tables which are used to determines the routing of data streams) for a second route (explained above in Aicklen) by identifying, in a consecutive fashion (sequential search is considered to be in a consecutive fashion) within the at least two timeslot interchanges (disclosed in Aicklen and Eyeson above), the p timeslots (timeslots T2 and T5 in Aicklen as explained above).

Regarding claim 9, performing a uniform search (see "a sequential table search method" recited in paragraph 72, where the tables are VPN identification tables which

are used to determines the routing of data streams) for a second route (explained above in Aicklen) by determining whether the p timeslots have less load than loads of remaining timeslots within the at least two timeslot interchanges (explained in Aicklen above).

Regarding claim 11, performing a sequential search (see “a sequential table search method” recited in paragraph 72, where the tables are VPN identification tables which are used to determines the routing of data streams) for a second route (explained above in Aicklen) by identifying, in a consecutive fashion (sequential search is considered to be in a consecutive fashion) within the at least two timeslot interchanges (disclosed in Aicklen and Eyeson above), the nth timeslot (timeslots T2 and T5 in Aicklen as explained above).

Regarding claim 12, performing a uniform search (see “a sequential table search method” recited in paragraph 72, where the tables are VPN identification tables which are used to determines the routing of data streams) for a second route (explained above in Aicklen) by determining whether the nth timeslot has less load than loads of remaining timeslots within the at least two timeslot interchanges (explained in Aicklen above).

Regarding claim 19, performing a sequential search (see “a sequential table search method” recited in paragraph 72, where the tables are VPN identification tables which are used to determines the routing of data streams) for the second route (explained above in Aicklen) by identifying, in a consecutive fashion (sequential search is considered to be in a consecutive fashion) within the at least two timeslot



interchanges (disclosed in Aicklen and Eyeson above), the p timeslots within the second route (timeslots T2 and T5 in Aicklen as explained above).

Regarding claim 20, performing a uniform search (see "a sequential table search method" recited in paragraph 72, where the tables are VPN identification tables which are used to determine the routing of data streams) for a second route (explained above in Aicklen) by determining whether the nth timeslot has less load than loads of remaining timeslots within the at least two timeslot interchanges (explained in Aicklen above).

It would have been obvious for one of the ordinary skill in the art at the time of the invention to modify the system of Williams using features, as taught by Aicklen, Eyeson and Akahane, in order to provide a non-blocking multi-rate transmission in a communication network involving different type of signals and prevent packets being dropped from failed or congested path.

10. Claim 16 is rejected under 35 U.S.C. 103(a) as being unpatentable over Williams (US 6,356,550) in view of Lappetelamen (US 6,693,915).

Williams discloses the claimed limitations above.

Williams also discloses the following features, searching one timeslot if the connection is of the first type (see Fig. 2J, data of the first type A is placed in timeslots 1).

Williams does not disclose the following features: regarding claim 16, determining whether a connection is one of the first type and a second type; and

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searching for one of (m-1) timeslots and (m-3) timeslots if the connection is of the first type.

Lappetelamen discloses a method of efficient bandwidth allocation for high-speed wireless data transmission system including the following features.

Regarding claim 16, determining (see "transmits information...on the connection type" recited in column 8, line 34-37) whether a connection is one of the first type and a second type (see "e.g. multimedia connection, data connection" recited in column 8, line 34-37); and searching for one of (m-1) timeslots and (m-3) slots if the connection is of the first type (see "The type of the connection...affect...the number of timeslots...to be allocated for the connection" recited in column 8, line 37-39).

Although, Lappetelamen does not specifically disclose using one of (m-1) timeslots and (m-3) timeslots for the connection of the first type, it is determined any number of timeslots can be used for a connection of the first type. Therefore, it would have been obvious to one of ordinary skill in the art at the time of the invention to modify the number of timeslots used for the connection of the first type in order to transmit the provide the required quality of service to the particular type of connection. The specific number of timeslots used recited in the claimed invention is not patently distinct from the reference of Lappetelamen because:

MPEP 2144.04, section IV A states:

#### "IV. CHANGES IN SIZE, SHAPE, OR SEQUENCE OF ADDING INGREDIENTS

##### A. Changes in Size/Proportion

In *Gardner v. TEC Systems, Inc.*, 725 F.2d 1338, 220 USPQ 777 (Fed. Cir. 1984), cert. denied, 469 U.S. 830, 225 USPQ 232 (1984), the Federal Circuit held that, where the only difference between the prior art and the claims was a recitation of relative dimensions of the claimed device and a device having the claimed relative dimensions would not perform differently than the prior art device, the claimed device was not patentably distinct from the prior art device."

It would have been obvious to one of the ordinary skill in the art at the time of the invention to modify the system of Williams using features, as taught by Lappetelamen, in order to provide the desired quality of service for the different type of connections.

11. Claim 21-25 are rejected under 35 U.S.C. 103(a) as being unpatentable over Williams (US 6,356,550) in view of Boily (US 2004/0001454).

Williams discloses a flexible time division multiplexed bus using SONET formatting including the following features.

Regarding claim 21, a system for switching multi-rate communications (see "communications system transports facility signals including: DS1, DS2, and DS3 signals, a 2.048 Mbps E1 signal...in addition to the 51.84 Mbps SONET STS-1 level signal" recited in column 6, line 63-67) comprising: at least a first data collection (see data collection X in Fig. 2G and 2J) having a common first type (see Fig. 2G, data collection 2G is of the common VT2 type), wherein the first data collection includes m sets of data (see Fig. 2G, where there are 4 sets of X data, that is,  $m = 4$ ); p sets of data of the first data collection are communicated to a first set of p timeslots (see Fig. 2J,

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wherein 3 sets of data collection X are filled into 3 timeslots, timeslots 2, 24, and 46; that is,  $p = 3$ ) and from which at least one overflow set of data from the first data collection is communicated to at least one overflow timeslot (see Fig. 2J, a fourth set of data X is filled in timeslot 68) wherein  $p$  is less than  $m$  ( $p = 3$ ,  $m = 4$  as explained above).

Regarding claim 22, a system for switching multi-rate communications (see "communications system transports facility signals including: DS1, DS2, and DS3 signals, a 2.048 Mbps E1 signal...in addition to the 51.84 Mbps SONET STS-1 level signal" recited in column 6, line 63-67) comprising: receive a plurality of connections (see the multiple connections, A-D, X-Z, M-O in Fig. 2F-2I) and transform the plurality of connections into a system data format (see Fig. 2J-2K, where the VT1, VT2, VT3, and VT6 connections shown in Fig. 2F-2I are transformed into a single SONET STS-1 format).

Regarding claim 24, receive at least a first, a second, and a third data collection (see data collections X, Y and Z in Fig. 2G and 2J), wherein each of the first, the second, and the third data collection have a common first type (see Fig. 2G, X, Y and Z are of the type VT2), and each of the first, the second, and the third data collection include  $m$  sets of data (see Fig. 2G, X, Y, and Z each include four sets of data); and  $p$  sets of data from the first data collection are communicated to a first set of  $p$  timeslots (see Fig. 2J-2K, timeslot 2, 24 and 46, each filled with data from data collection X,  $p = 3$ ),  $p$  sets of data from the second data collection are communicated to a second set of  $p$  timeslots (see Fig. 2J, timeslot 10, 32, and 53 are filled with three sets of data from

the data collection Y, which is shown in Fig. 2G), p sets of data from the third data collection are communicated to a third set of p timeslots (see Fig. 2J, timeslot 17, 37, and 61 are filled with three sets of data from the data collection Z, which is shown in Fig. 2G), an mth set of data from the first data collection (see Fig. 2G, the fourth set of data in collection X) are communicated to a fourth set of p timeslots (see Fig. 2J, timeslots 68, 77, and 82, where timeslot 68 contains the fourth set of data X), an mth set of data from the second data collection are communicated to the fourth set of p timeslots (see Fig. 2J, timeslots 68, 77, and 82, where timeslot 77 contains the fourth set of data Y), and an mth set of data from the third data collection are communicated to the fourth set of p timeslots (see Fig. 2J, timeslots 68, 77, and 82, where timeslot 82 contains the fourth set of data X).

Regarding claim 25, receive a plurality of connections (see the multiple connections, A-D, X-Z, M-O in Fig. 2F-2I) and transform the plurality of connections into a system data format (see Fig. 2J-2K, where the VT1, VT2, VT3, and VT6 connections shown in Fig. 2F-2I are transformed into a single SONET STS-1 format).

Williams does not disclose the following features: regarding claim 21, a time-space switch element configured to receive at least a first data collection; and a buffer from which p sets of data of the first data collection are communicated; regarding claim 22, an input interface is configured to be coupled to said time-space switch element, wherein said interface is configured to receive a plurality of connections and transform the plurality of connections into a system data format; an output interface configured to be coupled to said time-space switch element, said output interface configured to

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transform an output of said time-space switch element into the plurality of connections; regarding claim 23, a memory configured to store a program, wherein the program is configured to determine whether a timeslot interchange that includes the memory has a capacity to support a number of connections of the first type and a second type; regarding claim 24, a time-space switch element configured to receive the data collections; and a buffer from which the sets of data are to be communicated towards; regarding claim 22, an input interface is configured to be coupled to said time-space switch element, wherein said interface is configured to receive a plurality of connections and transform the plurality of connections into a system data format; an output interface configured to be coupled to said time-space switch element, said output interface configured to transform an output of said time-space switch element into the plurality of connections.

Boily discloses a timeslot interchange switch including the following features.

Regarding claim 21, a time-space switch element (see "timeslot interchange switch" recited in the abstract and shown in Fig. 1) configured to receive at least a first data collection (explained in Williams above, also see incoming stream 45 and ingress link 12 shown in Fig. 1 and 2); and a buffer (see ingress buffer 30 and egress buffer 32 shown in Fig. 3) from which p sets of data of the first data collection are communicated (explained in Williams above, also see incoming stream 45 and ingress link 12 shown in Fig. 1 and 2).

Regarding claim 22, an input interface (see ingress stage 24, ingress buffer 30 and cross connect 28 in Fig. 3, where the time-space switching occurs) is configured to

be coupled to said time-space switch element, wherein said interface is configured to receive a plurality of connections and transform the plurality of connections into a system data format (explained in above in Williams); an output interface (see Fig. 3, egress stage 26 including egress buffer 32) configured to be coupled to said time-space switch element (see Fig. 3, egress stage 26 is part of the TSI system), said output interface configured to transform an output of said time-space switch element into the plurality of connections (see Fig. 4, where the egress buffer 32A is shown, including active locations 48 that holds active connection as explained in paragraph 31-32; each column of egress buffer 32 A containing active locations 48 are output to an output 14 as explained in paragraph 28).

Regarding claim 23, a memory (see "field programmable gate array (FPGA)" recited in paragraph 38) configured to store a program (see System 10 may be implemented a...FPGA...The invention may be embodied in a set of instructions for configuring an FPGA...to provide apparatus according to the invention" recited in paragraph 38), wherein the program is configured to determine whether a timeslot interchange (see Fig. 1-2, both shows a timeslot interchange switch) has a capacity to support a number of connections (see "in cases where there is an upper limit to the number of connections...system 10 may maintain a count of active locations" recited in paragraph 43) of the first type and a second type (see "switching data of various types..." recited in paragraph 53).

Regarding claim 24, a time-space switch element (see "timeslot interchange switch" recited in the abstract and shown in Fig. 1) configured to receive data collections

(explained in Williams above, also see incoming stream 45 and ingress link 12 shown in Fig. 1 and 2); and a buffer (see ingress buffer 30 and egress buffer 32 shown in Fig. 3) from which the sets of data are to be communicated towards (explained in Williams above, also see incoming stream 45 and ingress link 12 shown in Fig. 1 and 2).

Regarding claim 25, an input interface (see ingress stage 24, ingress buffer 30 and cross connect 28 in Fig. 3, where the time-space switching occurs) is configured to be coupled to said time-space switch element, wherein said interface is configured to receive a plurality of connections and transform the plurality of connections into a system data format (explained in above in Williams); an output interface (see Fig. 3, egress stage 26 including egress buffer 32) configured to be coupled to said time-space switch element (see Fig. 3, egress stage 26 is part of the TSI system), said output interface configured to transform an output of said time-space switch element into the plurality of connections (see Fig. 4, where the egress buffer 32A is shown, including active locations 48 that holds active connection as explained in paragraph 31-32; each column of egress buffer 32 A containing active locations 48 are output to an output 14 as explained in paragraph 28).

It would have been obvious for one of the ordinary skill in the art at the time of the invention to modify the system of Williams using features, as taught by Boily, in order to provide a non-blocking multi-rate transmission in a communication network involving different type of signals.



***Allowable Subject Matter***

12. Claim 13-15 are objected to as being dependent upon a rejected base claim, but would be allowable if rewritten in independent form including all of the limitations of the base claim and any intervening claims and rewritten to overcome the rejection(s) under 35 U.S.C. 112, 2nd paragraph, set forth in this Office action and to include all of the limitations of the base claim and any intervening claims.

***Conclusion***

13. The prior art made of record and not relied upon is considered pertinent to applicant's disclosure.

Dally (US 2001/0033569) discloses a multistage digital cross connect with integral frame timing.

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Ju-Tai Kao whose telephone number is (571)272-9719. The examiner can normally be reached on Monday ~Friday 7:30 AM ~5:00 PM EST.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Kwang Yao can be reached on (571)272-3182. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

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Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

Ju-Tai Kao



KWANG BIN YAO  
SUPERVISORY PATENT EXAMINER

